Lecture 22: Cascades Framework

CREATING THE NEXT®

Today's Agenda

Recap

Logical Query Optimization

Physical Query Optimization

Cascades Optimizer

Case Studies

Conclusion





Optimization Search Strategies

- Choice 1: Heuristics
 - INGRES, Oracle (until mid 1990s)
- Choice 2: Heuristics + Cost-based Join Search
 - ▶ System R, early IBM DB2, most open-source DBMSs
- Choice 3: Randomized Search
 - ► Academics in the 1980s, current Postgres
- Choice 4: Stratified Search
 - ► IBM's STARBURST (late 1980s), now IBM DB2 + Oracle
- Choice 5: Unified Search
 - Volcano/Cascades in 1990s, now MSSQL + Greenplum



Stratified Search

- First rewrite the logical query plan using transformation rules.
 - ▶ The engine checks whether the transformation is allowed before it can be applied.
 - Cost is never considered in this step.
- Then perform a cost-based search to map the logical plan to a physical plan.



- Unify the notion of both logical \rightarrow logical \rightarrow physical transformations.
 - ▶ No need for separate stages because everything is transformations.
- This approach generates a lot more transformations so it makes heavy use of memoization to reduce redundant work.



Top-Down vs. Bottom-Up

• Top-down Optimization

- Start with the final outcome that you want, and then work down the tree to find the optimal plan that gets you to that goal.
- ► Example: Volcano, Cascades

Bottom-up Optimization

- Start with nothing and then build up the plan to get to the final outcome that you want.
- Examples: System R, Starburst



Logical Query Optimization

- Transform a logical plan into an equivalent logical plan using pattern matching rules.
- The goal is to increase the likelihood of enumerating the optimal plan in the search.
- Cannot compare plans because there is no cost model but can "direct" a transformation to a preferred side.



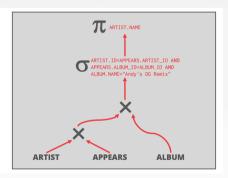
- Split Conjunctive Predicates
- Predicate Pushdown
- Replace Cartesian Products with Joins
- · Projection Pushdown
- Reference



Split Conjunctive Predicates

SELECT ARTIST.NAME
FROM ARTIST, APPEARS, ALBUM
WHERE ARTIST.ID=APPEARS.ARTIST_ID
AND APPEARS.ALBUM_ID=ALBUM.ID
AND ALBUM.NAME="Andy's OG Remix"

Decompose predicates into their simplest forms to make it easier for the optimizer to move them around.

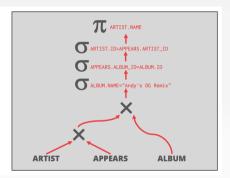




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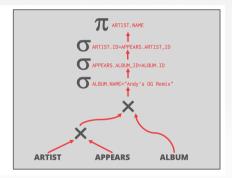




Predicate Pushdown

SELECT ARTIST.NAME FROM ARTIST, APPEARS, ALBUM WHERE ARTIST.ID=APPEARS.ARTIST_ID AND APPEARS.ALBUM_ID=ALBUM.ID AND ALBUM.NAME="Andy's OG Remix"

Move the predicate to the lowest point in the plan after Cartesian products.

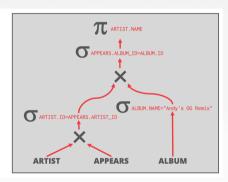




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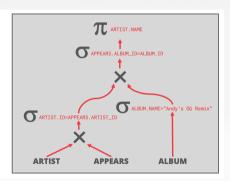




Replace Cartesian Products with Joins

SELECT ARTIST.NAME FROM ARTIST, APPEARS, ALBUM WHERE ARTIST. ID=APPEARS. ARTIST ID AND APPEARS.ALBUM_ID=ALBUM.ID AND ALBUM NAME="Andy's OG Remix"

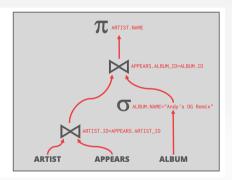
Replace all Cartesian Products with inner joins using the join predicates.





SELECT ARTIST.NAME
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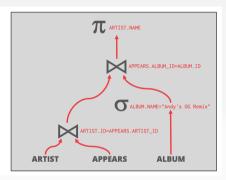




Projection Pushdown

SELECT ARTIST.NAME
FROM ARTIST, APPEARS, ALBUM
WHERE ARTIST.ID=APPEARS.ARTIST_ID
AND APPEARS.ALBUM_ID=ALBUM.ID
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Eliminate redundant attributes before pipeline breakers to reduce materialization cost.

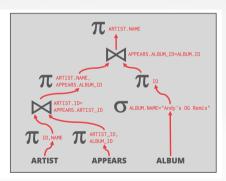




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Physical Query Optimization

- Transform a query plan's logical operators into physical operators.
 - Add more execution information
 - Select indexes / access paths
 - Choose operator implementations
 - ► Choose when to materialize (*i.e.*, temp tables).
- This stage must support **cost model estimates**.



Observation

- All the queries we have looked at so far have had the following properties:
 - Equi/Inner Joins
 - Simple join predicates that reference only two tables.
 - No cross products
- Real-world queries are much more complex:
 - Outer Joins
 - Semi-joins
 - Anti-joins



- No valid reordering is possible.
- The $A \supset B$ operator is not commutative with $B \supset C$.
 - ► The DBMS does not know the value of B.val (may be *NULL*) until after computing the join with A.
- Reference

```
SELECT * FROM
A LEFT OUTER JOIN B
ON A.id = B.id
FULL OUTER JOIN C
ON B.val = C.id;
```



Plan Enumeration

· Approach 1: Transformation

Modify some part of an existing query plan to transform it into an alternative plan that is equivalent.

• Approach 2: Generative

- Assemble building blocks to generate a query plan (similar to dynamic programming).
- Reference



Dynamic Programming Optimizer

- Model the query as a hypergraph and then incrementally expand to enumerate new plans.
- Algorithm Overview:
 - ► Iterate connected sub-graphs and incrementally add new edges to other nodes to complete query plan.
 - ▶ Use rules to determine which nodes the traversal is allowed to visit and expand.
- Reference



Cascades Optimizer

- Object-oriented implementation of the Volcano query optimizer.
- Materialize transformations on the fly (rather than pre-generate them all at once).
- Unlike Volcano, restricts the set of transformations to constrain the search space.
- Supports simplistic expression re-writing through a direct mapping function rather than an **exhaustive search**.



Cascades Optimizer

- Optimization tasks as data structures.
- Rules to place **property enforcers** (*e.g.*, sorting order).
- Ordering of transformations by priority. Dynamically adjust ordering as we traverse the search tree.
- Predicates are first class citizens (same as logical/physical operators).



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- An **expression** is an **operator** with zero or more input expressions.
- Logical Expression: $(A \bowtie B) \bowtie C$
- **Physical Expression**: $(A_{Seq} \bowtie_{HJ} B_{Seq}) \bowtie_{NL} C_{Idx}$

SELECT * FROM A

JOIN B ON A.id = B.id

JOIN C ON C.id = A.id;



Cascades - Groups

- A **group** is a set of **logically equivalent** logical and physical expressions that produce the same output.
 - All logical forms of an expression.
 - All physical expressions that can be derived from selecting the allowable physical operators for the corresponding logical forms.





Cascades - Multi-Expression

- Instead of explicitly instantiating all possible expressions in a group, the
 optimizer implicitly represents redundant expressions in a group as a
 multi-expression.
 - This reduces the number of transformations, storage overhead, and repeated cost estimations.
 - ▶ We can make decisions about whether to traverse [AB] first vs. [C] first.

```
| Logical Multi-Exps | Physical Multi-Exps | 1. [AB] ⋈ [C] | 1. [AB] ⋈ [C] | 2. [AB] ⋈ [C] | 2. [AB] ⋈ [C] | 3. [AB] ⋈ [C] | 4. [A] ⋈ [BC] | 4. [A] ⋈ [BC] | 4. [BC] ⋈ [A] ⋈ [A] | 1. [A] ⋈ [AB] ⋈ [A
```

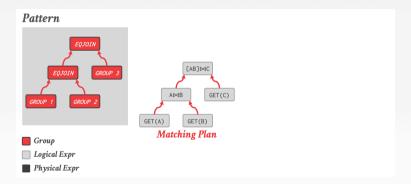


Cascades - Rules

- A **rule** is a transformation of an expression to a logically equivalent expression.
 - ► Transformation Rule: Logical to Logical
 - ► Implementation Rule: Logical to Physical
- Each rule is represented as a pair of attributes:
 - Pattern: Defines the structure of the logical expression that can be applied to the rule.
 - **Substitute:** Defines the structure of the result after applying the rule.

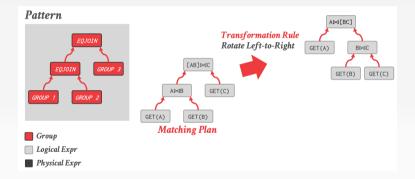


Cascades - Rules



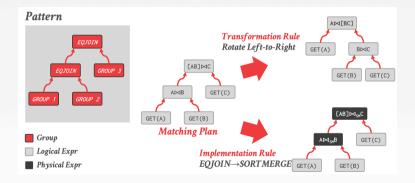


Cascades - Rules





Cascades – Rules





Cascades - Memo Table

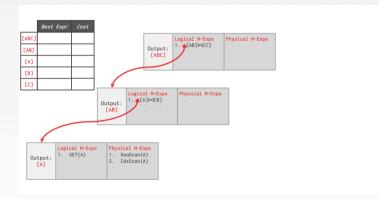
- Stores all previously explored alternatives in a compact graph structure / hash table.
- Equivalent operator trees and their corresponding plans are stored together in groups.
- Provides memoization, duplicate detection, and property + cost management.



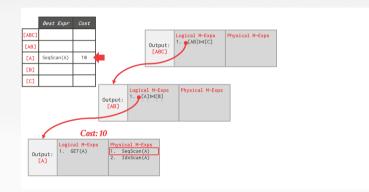
Principle of Optimality

- Every sub-plan of an optimal plan is itself optimal.
- This allows the optimizer to restrict the search space to a smaller set of expressions.
 - ► The optimizer never has to consider a plan containing sub-plan *P*1 that has a greater cost than equivalent plan *P*2 with the same physical properties.
 - Reference

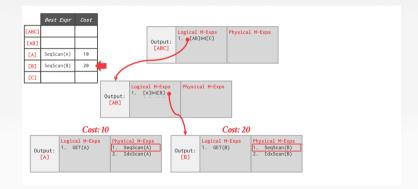




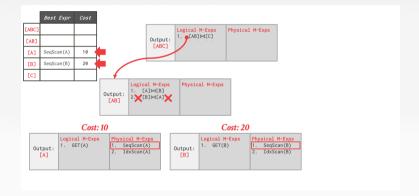




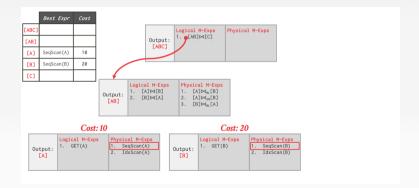




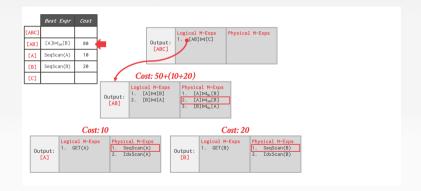




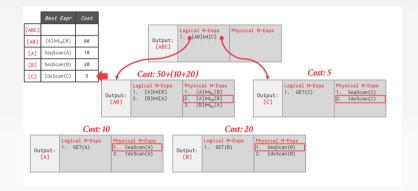




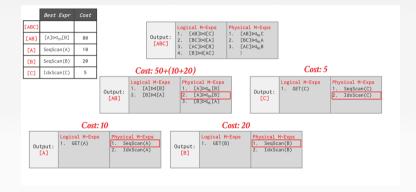














Search Termination

- Approach 1: Wall-clock Time
 - Stop after the optimizer runs for some length of time.
- Approach 2: Cost Threshold
 - ► Stop when the optimizer finds a plan that has a lower cost than some threshold.
- Approach 3: Transformation Exhaustion
 - Stop when there are no more ways to transform the target plan. Usually done per group.



Case Studies

Cascades: Implementations

• Approach 1: Standalone Optimizer Generator

- ► Wisconsin OPT++ (1990s)
- Portland State Columbia (1990s)
- ► Pivotal Orca (2010s)
- Apache Calcite (2010s)

Approach 2: Integrated

- ► Microsoft SQL Server (1990s)
- ► Tandem NonStop SQL (1990s)
- ► Clustrix (2000s)
- ► CMU Peloton (2010s RIP)



Pivotal Orca

- Standalone Cascades (Optimization-as-a-service).
- Reference
 - Originally written for Greenplum.
 - Extended to support HAWQ.
- A DBMS can use Orca by implementing API to send catalog + stats + logical plans and then retrieve physical plans.
- Supports multi-threaded search.



Orca - Engineering

• Issue 1: Remote Debugging

- Automatically dump the state of the optimizer (with inputs) whenever an error occurs.
- ► The dump is enough to put the optimizer back in the exact same state later for further debugging.

Issue 2: Optimizer Accuracy

 Automatically check whether the ordering of the estimate cost of two plans matches their actual execution cost.



Apache Calcite

- Standalone extensible query optimization framework for data processing systems.
 - Support for pluggable query languages, cost models, and rules.
 - Does not distinguish between logical and physical operators. Physical properties are provided as annotations.
- Reference
- Originally part of LucidDB.



• Rewriter

Logical-to-logical transformations with access to the cost-model.

Enumerator

- Logical-to-physical transformations.
- Mostly join ordering.

Planner

- Convert physical plans back to SQL.
- Contains MemSQL-specific commands for moving data.
- Reference



Conclusion

Parting Thoughts

- Cascades
 - Optimization tasks as data structures.
 - ► Rules to place **property enforcers** (*e.g.*, sorting order).
 - Ordering of transformations by priority.
 - Predicates are first class citizens (same as logical/physical operators).
- All of this relies on a good **cost model**.
- A good cost model needs good statistics.



Next Class

• Non-Traditional Query Optimization Techniques

