High Performance Computing: Tools and Applications

Edmond Chow School of Computational Science and Engineering Georgia Institute of Technology

Lecture 10

Don't forget about:

KMP\_AFFINITY=verbose,none
KMP\_AFFINITY=verbose,compact
KMP\_AFFINITY=verbose,granularity=fine,compact



# Intel Xeon Phi



Logically an UMA architecture.

OpenMP is good at parallelizing loops.

What if we want to parallelize

- while loops where we don't know the number of iterations in advance
- traversing linked lists
- recursive functions
- batch of tasks, where performing a task can create new tasks

- OpenMP has the facility to define a pool of tasks
- These tasks are executed by threads when threads are free
- Tasks can be added to the pool dynamically

#### First example

```
#pragma omp parallel
#pragma omp single
 printf("Start creating tasks by thread %d\n", omp_get_thread_num());
 #pragma omp task
 printf("1st task executed by thread %d\n", omp_get_thread_num());
 #pragma omp task
 printf("2nd task executed by thread %d\n", omp_get_thread_num());
 #pragma omp task
 printf("3rd task executed by thread %d\n", omp_get_thread_num());
 printf("Done creating tasks\n");
 // taskwait is needed if a barrier is desired
 #pragma omp taskwait
 printf("All tasks completed by now.\n");
```

- the task may be run by a thread immediately or added to pool of tasks to be executed later (depends on runtime system)
- tasks are not consumed in any order
- when a task is created, it takes up resources including memory for private variables
- some omp task clauses
  - ▶ if
  - untied
  - ▶ default
  - ▶ mergeable
  - ▶ private
  - firstprivate
  - shared

- if a variable is private in the enclosing context, then the default is that it is firstprivate in the task
- ► if task clause is private, then initial copy is not performed

#### Second example – there is a bug in this code

```
int taskid;
#pragma omp parallel
#pragma omp single
  for (int i=0; i<10; i++) {</pre>
    taskid = i+1;
    #pragma omp task
    printf("thread %2d: task %2d\n",
      omp_get_thread_num(), taskid);
  printf("Done creating tasks\n");
// implied barrier
int threadid = omp_get_thread_num();
if (threadid == 0) printf("All tasks completed by now.\n");
```

```
#pragma omp parallel
#pragma omp single
  for (int i=0; i<10; i++) {
   int taskid = i+1; // taskid is private
    #pragma omp task // for each task, taskid is threadprivate
   printf("thread %2d: task %2d\n",
     omp_get_thread_num(), taskid);
 printf("Done creating tasks\n");
// implied barrier
int threadid = omp_get_thread_num();
if (threadid == 0) printf("All tasks completed by now.\n");
```

### Fibonnaci numbers (recursive algorithm)

```
// 1 1 2 3 5 8 13 ...
int fib(int n)
    int x, y;
    if (n < 2) return n;
    x = fib(n-1);
    y = fib(n-2);
    return x+y;
void main()
    int n = 20;
    printf("fib(%d) = %d\n", n, fib(n));
```

```
int fib(int n) {
    int x, y;
    if (n < 2) return n;</pre>
#pragma omp task shared(x)
    x = fib(n-1);
#pragma omp task shared(y)
    y = fib(n-2);
#pragma omp taskwait
    return x+y;
void main() {
    int n = 20;
#pragma omp parallel
#pragma omp single
    printf("fib(%d) = %d\n", n, fib(n));
```

Note: without shared clause, x and y would be firstprivate in the tasks because x and y are private in the enclosing context.

# Cilk Plus features of Intel compilers

- Cilk Plus defines three keywords
  - ▶ cilk\_spawn
  - cilk\_sync
  - ▶ cilk\_for
- Philosophy is for the programmer to expose parallelism, and let the runtime decide how to optimize thread scheduling, vectorization, etc.
- Cilk Plus is arguably more task-based than OpenMP. Task-stealing, which is transparent to the user, is used for balancing load.
- Cilk Plus also defines array notation (facilitates both multithreading and vectorization) and reducers (parallel data types that help avoid the use of locks)
- On Intel compilers, simply include Cilk Plus header files, generally cilk/cilk.h for nicer keywords. No special compilation flags are needed

```
int fib(int n)
   int x, y;
    if (n < 2) return n;</pre>
   x = cilk_spawn fib(n-1);
   y = fib(n-2); // No cilk_spawn needed here.
   cilk_sync; // Block here until all
                 // spawned functions are complete.
   return x+y;
    // Implied cilk_sync at end of any function
    // that contains cilk_spawn.
void main()
   int n = 20;
   printf("fib(%d) = %d\n", n, fib(n));
```

# What happens with cilk\_spawn

- Only functions can be spawn, not sections of code. This makes it easier (than OpenMP) to define the data environemnt (e.g., no need for threadprivate). This also makes it unnecessary to use pragmas to define sections of code.
- Each thread has its own work queue
- When a thread encounters cilk\_spawn, then a task, which is the continuation of the original task is added to the end of its own work queue; the thread executes the code that is spawn
- When a thread has finished a task, it takes a new task from the end of the queue (this is better for cache usage)
- When threads do not have work, they steal tasks from other threads from the front of their queues
- In this sense, Cilk Plus does not force parallel execution, but provides the opportunity for parallel execution
- Like OpenMP, removing the keywords should give a correct serial program

```
int i;
cilk_for (i=0; i<10; i++)
{
    int id = __cilkrts_get_worker_number();
    printf("iteration %2d: worker %2d\n", i, id);
}</pre>
```

- Using the worker number is discouraged in Cilk Plus
- Programmer should not worry about how threads are scheduled (trade-off between easier programming and performance)

- Divide and conquer
- The thread encountering does the following:
  - makes a task that is the first half of the iterations; adds task to its queue
  - for the remaining iterations, repeat the above, by making a task that is the first half of the remaining iterations, and adding this task to the queue
  - etc.
- Free threads will steal tasks from the front of another task's queue
- How much to steal?
  - stealing half of the total work is good for balancing load
  - this means stealing one item (more efficient than stealing many items)

- To sum an array of numbers with multiple threads, locks or similar mechanisms are needed
- Reducers in Cilk Plus are C++ classes that perform reduction operations without the need for the programmer to use locks or critical sections

```
// Sum the numbers 1-1000 in parallel,
// adding a pause to allow the continuation to be stolen
cilk::reducer< cilk::op_add<int> > parallel_sum(0);
cilk_for(int i = 0; i < 1000; i++)
{
    if (0 == i % 10)
        stall();
    *parallel_sum += i;
}
printf("Parallel sum: %d\n", parallel_sum.get_value());
```

Reference: www.cilkplus.org

# How is reduction implemented in Cilk Plus?

- Reduction is performed between two threads when they join.
- Thus the reduction is performed like a binary tree.

# Cilk Plus array notation

Array notation: A[start:length:stride]

Array notation implies independent operations and vector code will be generated

```
A[:] = 5; // set all elements of static array
A[7:3] = 4; // set elements 7, 8, 9
A[1:5:2] = 4; // set elements 1, 3, 5, 7, 9
A[:] = B[:] * C[:] + 5; // Cilk assumes no overlap
C[X][:] = A[:]; // X \text{ is an expression}
C[:] = A[B[:]]; // gather
A[B[:]] = C[:]; // scatter
C[:][:] = 12; // two-dimensional arrays
func(A[:]); // pass elements one-by-one
__sec_reduce_add(A[:]) // returns scalar
```

Reference: www.cilkplus.org

# Cilk Plus array notation

```
int a[array_size];
const char *results[array_size];
if (5 == a[:])
    results[:] = "Matched";
else
    results[:] = "Not Matched";
```

#### is equivalent to

```
int a[array_size];
const char *results[array_size];
for (int i = 0; i < array_size; i++)
{
    if (5 == a[i])
        results[i] = "Matched";
    else
        results[i] = "Not Matched";
}</pre>
```

Reference: www.cilkplus.org

- environment variable: CILK\_NWORKERS
- at run time: \_\_\_\_cilkrts\_set\_param("nworkers", "N")

Class on Tuesday, Sept. 27 is cancelled.